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**Improved upper and lower time bounds for parallel random access machines without simultaneous writes.**

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A PRAM or parallel random access machine consists of an infinite number of processors that operate in parallel and share an infinite memory. At each time step each processor reads a value from the shared memory, changes to a new state, and writes this new state back into the shared memory. Simultaneous read operations of a single shared-memory cell by many processors are allowed but simultaneous write operations into a single cell are not. An input to a PRAM  $M$  consists of a finite sequence  $I = (x_0, x_1, \dots, x_{n-1})$  of natural numbers, of which  $x_i$  is placed in cell  $i$  of the shared memory at the start of the computation of  $M$ . When the processors have terminated, the contents of cell 0 constitute the output computed by  $M$ . The time required by  $M$  is the maximum over all inputs of size  $n$  of the number of steps needed to process that input, considered as a function of  $n$ .

Using a communication argument, the authors improve by a constant factor known lower and upper bounds on the running time for PRAMs. More precisely, they show that any PRAM that computes a so-called critical function—i.e., a Boolean function for which there exists an input  $I$  with the property that changing a single bit  $x_i$  from  $I$  changes its output—requires at least time  $0.5 \log_2 n - O(1)$ , and that there exists a critical function which can be computed in time  $0.57 \log_2 n - O(1)$ .

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